Handling C++ Exceptions in MPI Applications

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1 INTRODUCTION AND MOTIVATION

Handling error states in C++ applications is managed by exceptions [1]. This technique protects a piece of code by a try-catch code block. Any error arising in the try block is converted into an exception object and resolved in the catch block. If the error is not handled, the application is still properly terminated. Exceptions bring many significant benefits over standard error code checking implemented in the C language, the most important of which is decoupled program logic and error handling.

In distributed MPI applications, it is often necessary to inform the other processes (ranks), that something wrong happened, and that the application should either recover from the faulty state, or report the error and terminate gracefully. Unfortunately, the MPI standard does not provide any support for distributed error handling. It only defines return codes for MPI calls. The distribution of the error message is then left to the developers. If they are not willing do so after every MPI call, the MPI runtime may be informed that all errors are fatal. This, however, leads to application termination and the error reporting from all ranks, which is not user friendly while running over thousands of ranks.

After dropping the C++ interface in the MPI-3.0 standard [8], the MPI exception interface for reporting the MPI errors was removed from the default branch as well, yet it can still be enabled during MPI compilation. Nevertheless, the MPI exceptions only cover the MPI routines and report to the local rank only. This brings a lot of troubles since any exception thrown by either the MPI runtime or the user code breaks the standard code path and jumps into the exception handler. This may cause several communication routines to be skipped leading to a deadlock. This is even an worse situation than an uncoordinated crash, since the application hangs and keeps consuming computing resources until manually terminated.

Surprisingly, there is very little work published on the error handling inside MPI applications. Most research concerns on solving MPI faulty states such as a node crash, lost interconnection, or detection and recovery of deadlock states [6, 7, 9]. However, what if the MPI itself is working properly but the application encountered an unexpected situation? As far as the author knows, there is only a

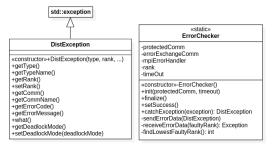


Figure 1: Class diagram of the MPI error checker.

single paper devoted to exception propagation in MPI applications by Engwer at.al [2]. This work uses a side channel for distributing error states over all ranks and asynchronous error checking. The technique seems to be pretty robust, but has a major drawback. All MPI calls have to be made via a custom interface which makes it unfit for applications using third party libraries.

This poster presents a new approach for exceptions handling in MPI applications. The goals are to (1) report any faulty state to the user in a nicely formatted way by just a single rank, (2) ensure the application will never deadlock, (3)propose a simple interface and ensure interoperability with other C/C++ libraries.

2 PROPOSED METHOD

The proposed method adopts a minimalistic interface. In the simplest case, the user can use only a single try-catch block to manage error reporting in a sensible way. Nevertheless, the user is free to use as many try-catch blocks as necessary.

The interface consists of two classes, DistException and Error-Checker, see Fig. 1. The DistException class wraps all exceptions derived from the base C++ class std::exception, including MPI::Exception and maintains information about the type of the exception, faulty MPI communicator, its name, faulty rank, user error code and error message, and the presence of deadlock.

The ErrorChecker class is used to report and handle exceptions. The application errors are reported by instances of DistException. In a nutshell, ErrorChecker clones the MPI communicator it is supposed to protect (e.g., MPI_COMM_WORLD) to build a side channel where the ranks announce to have passed a given checkpoint, usually at the end of the try block, at end of the iteration, or after a long operation the user should be informed about. At these points, the setSuccess() method is called to inform the others that no exception was thrown by this rank. If any other rank yet threw an exception, this routine will propagate this information within the local process by throwing a private instance of DistException.

To handle exceptions in the catch block, checkException() shall be called at the first place to pass on the information about

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local exceptions as well as to collect the error information from the other ranks, and to vote a suitable rank to log the error message (usually the root rank or the first alive). If the code is not deadlocked, the user can try to recover, or terminate the code gracefully otherwise. A code snippet showing the use can be seen in Listing 1.

int main(int argc, char** argv) 2 { // Init MPI and set the error handler. 3 MPI_Init(&argc, &argv); // Initialize error checker (static class) ErrorChecker::init(MPI_COMM_WORLD, timeout); // Protected block of the code 10 trv 11 { // Any combination of local and MPI computation may appear here. 12 13 MPI_Bast(...); foo():14 15 MPI_Barrier(...); 16 17 // The very last command of the try block set a success flag. 18 ErrorChecker::setSuccess(); 19 }// end of trv 20 // Error handling 21 catch (const std::exception& e) 22 23 // Find out whether any remote rank caused an exception and 24 25 // which rank is supposed to print out an err const auto& distException = ErrorChecker::catchException(e); 26 // Check whether code has deadlocked due to some blocking call 27 28 // or collective communication in progress. If so, find the rank // which will report the error, otherwise leave if for root. 29 const int reportingRank = (distException.getDeadlockMode() 30 ? distException.getRank()) 31 0; 32 if (reportingRank == myRank) reportError(distException); 33 // Print out error message and terminate or recover. 34 35 printErrorAndTerminate(distException); 36 }// end of catch 37 FrrorChecker::finalize(): 38 terminateApplication(EXIT_SUCCESS); 39 40 }// end of main

Listing 1: Simple usage of the error checking code

3 IMPLEMENTATION DETAILS

As mentioned previously, ErrorChecker creates a side channel to distribute information about exceptions.

The setSuccess() method informs all other ranks that no exception was thrown by this rank up to this point. For this purpose, a non-blocking MPI_Iallreduce call with an MPI_MIN operation is used. Each healthy rank sends the MAX_INT value to announce no error. Faulty ranks will, however, skip the setSuccess() method, but join the communication in the catchException() method and provide its rank instead. Back in setSuccess(), all ranks are sitting in MPI_Wait, waiting for the result of the reduce operation. If the result is MAX_INT, no error happened. Otherwise, the ranks will get the information about the first faulty rank that threw an exception. Although it would be possible to collect all faulty ranks by MPI_Iallgather, this is usually not required for reporting purposes. Once the faulty rank is known, the reporting rank asks the faulty one about the exception details using a set of pointto-point nonblocking communications. The exception details are only stored in the reporting rank, the others only know an error happened. This behaviour can be simply changed to an inform-all mode. Once all healthy ranks know about the remote exception,

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they create a local instance of DistException and throw it at the end of setSuccess() to join the faulty ranks in the catch block.

The catchException(const std::exception& e) first checks the exception type. If the exception was thrown by setSuccess() method, it was obviously created by a remote rank and rethrown locally. In that case, the exception details are extracted and returned for local processing. Otherwise, it must be a new local exception thrown somewhere in the try block and has to be propagated to the other ranks. This is done by joining the already mentioned MPI_Iallreduce communication. Since the faulty rank left the predescribed code path, it is possible that the other ranks will never join this call (never call setSuccess()) since being stuck in another blocking call. This is solved by periodic testing of the MPI_Iallreduce communication in catchException(). If not finished within a predefined timeout (10s by default), the application has deadlocked and the only solution is to call MPI_Abort. However, before that, it is necessary to report the error message and create an error code. Since there might be multiple faulty ranks (e.g., input file is not accessible from one node), it is necessary to find all faulty, but still alive ranks, to agree which one will overtake the reporting and termination responsibility. This is done by probing all ranks within the protected communicator. Each faulty rank sends its rank value to all other ranks and waits for the answer. Simply said, a custom deadlock-free variant of MPI_Iallgather communication is performed. After a given timeout, communications with not responding ranks are cancelled, and the reporter is chosen to be the process with the lowest rank. Consequently an exception information is wrapped into an MPI message with the deadlock mode set to true and sent to the reporter.

If deadlock was not detected while running MPI_Iallreduce, the faulty rank with the lowest rank value wraps the exception into the error message and sends it to the root rank. In this situation, the code is able to recover.

4 CONCLUSIONS

The proposed exception handling mechanism was integrated into the MPI version of the acoustics k-Wave toolbox [5]. The code was tested under different MPI implementations such as IntelMPI 19.x and OpenMPI 4.x up to 1536 ranks. As external libraries heavily utilising collective communications, distributed version of the fast Fourier transform (FFTW) [4] and the HDF5 [3] I/O libraries were chosen. The code was tested with several injected errors into multiple ranks such as non existing input file, disk quota exceeded, wrong rank in the MPI call, custom k-Wave exceptions, and standard system exceptions such as out of memory problems, numerical errors, etc. In all situations the code has worked properly.

The advantages of the proposed solution is that no dedicated rank for testing the errors is necessary, a single reduce operation is only required to confirm the application passed a check point, deadlock in application cannot interrupt the error handling, and the application always terminates gracefully with a proper error message. The necessary support for MPI exceptions to handle MPI error states, not part of the default branch, may be seen as a disadvantage, but it can be overcome by custom MPI error handlers. The code can be downloaded from https://github.com/jarosjir/MPIErrorChecker Handling C++ Exceptions in MPI Applications

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